

Fundamentals of Machine Learning

Yassine Laguel

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Maximal Inequality for bounded losses

Proposition : Upperbounding the max of bounded r.v.

Let $X_1, \dots, X_n$ be n centered random variables ($\mathbb{E}[X_i] = 0$ for all i), such that $X_i \in [a, b]$ almost surely for all $i \in \{1, \dots, n\}$.

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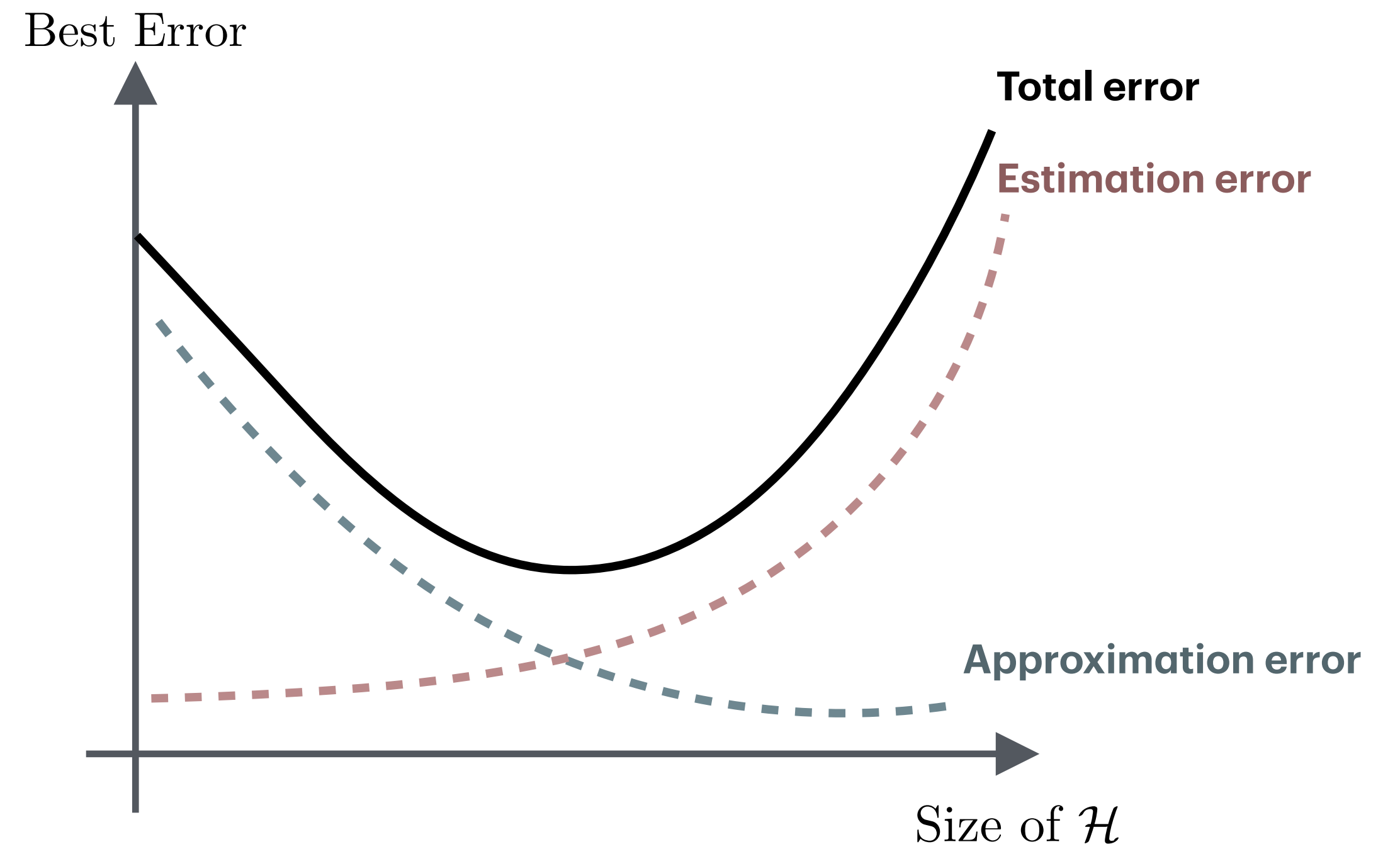
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- When $|\mathcal{H}|$ is infinite, this bound is not exploitable.

1. Supervised Learning Setting
2. Estimation vs Approximation
3. Maximal inequalities
4. Rademacher Complexity



Rademacher Complexity of a set

- Rademacher variables and complexity of a set

Rademacher Complexity of a set

■ Rademacher variables and complexity of a set

Definition : Rademacher random variables

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The Rademacher complexity of a set $T \subset \mathbb{R}^n$, denoted $\text{Rad}(T)$, is defined as:

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Proposition : Homotetie and translation

Let $T \subset \mathbb{R}^n, v \in \mathbb{R}^n, c \in \mathbb{R}$, and define $cT + v = \{ct + v, t \in T\}$.

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For any sets $T, T' \subset \mathbb{R}^d$, we have

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Proposition : Talagrand's lemma

Let $f_1, \dots, f_n : \mathbb{R} \rightarrow \mathbb{R}$ be n G -lipchitz functions.

Let $T \subset \mathbb{R}^n$ and define $(f_1, \dots, f_n) \circ T$ as

$$(f_1, \dots, f_n) \circ T = \{(f_1(t_1), \dots, f_n(t_n)), t \in T\}.$$

Then,

$$\text{Rad}((f_1, \dots, f_n) \circ T) \leq G \text{Rad}(T).$$

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Let $T \subset \mathbb{R}^n, v \in \mathbb{R}^n, c \in \mathbb{R}$, and define $cT + v = \{ct + v, t \in T\}$.

Then,

$$\text{Rad}(\text{conv}(T)) = \text{Rad}(T).$$

Proof : on the black board.

Proposition : Massart's Lemma - Complexity for finite Sets

Let $T \subset \mathbb{R}^n$ be finite.

Then,

$$\text{Rad}(T) \leq \max_{t \in T} \|t\|_2 \frac{\sqrt{2 \log(|T|)}}{n}.$$

Proof : on the black board.

Remark :

- The two previous propositions allow to upper bound the Rademacher complexity of the containing polyhedra.



Proposition : Talagrand's lemma

Let $f_1, \dots, f_n : \mathbb{R} \rightarrow \mathbb{R}$ be n G -lipchitz functions.

Let $T \subset \mathbb{R}^n$ and define $(f_1, \dots, f_n) \circ T$ as

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Remark :

- This lemma will later be used to establish statistical bounds when the loss ℓ satisfies a Lipschitz condition.

Rademacher complexity of a hypothesis class

■ Formal definition & Symmetrization Lemma

Definition : Rademacher complexity of a class of functions

Let $(\Omega_1, \dots, \Omega_n)$ be n i.i.d. Rademacher variables.

We define the empirical Rademacher complexity of the hypothesis class $\mathcal{H}$, denoted $\widetilde{\text{Rad}}(\mathcal{H})$, as

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Under the Empirical Risk Minimization Framework, we have :

$$\mathbb{E} \left[\sup_{f \in \mathcal{H}} \{r(f) - R(f)\} \right] \leq 2 \text{Rad}(\mathcal{H}).$$

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■ Applications in the regression setting

Proposition : Complexity for Lipschitz Losses

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Proof : on the black board.

Remarks :

- Hence, assuming assuming a Lipschitz error function allows us to focus on the complexity spanned by the set of estimators rather than the composition $\ell(f(X), Y)$.

Linear Regression Examples

Proposition : Linear prediction with ℓ_2 -constraints, and ℓ_2 bounded data.

Let $\mathcal{H} \triangleq \{x \mapsto \omega^\top x, \|\omega\|_2 \leq c\}$, for $c > 0$ fixed.

Let $x = (x_1, \dots, x_n) \in \mathcal{X}^n$ be fixed and $T_x \triangleq \{(f(x_1), \dots, f(x_n)), f \in \mathcal{H}\}$

Then,

$$\text{Rad}(T_x) \leq c \cdot \frac{\max_{1 \leq i \leq n} \|x_i\|_2}{\sqrt{n}}.$$

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- We directly deduce

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Remark :

- Note the logarithmic dependency in the dimension!

Linear Regression Examples

Proposition : Linear prediction with ℓ_2 -constraints, and ℓ_2 bounded data.

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Proof : Almost identical to the previous one.

A Nonlinear Example - Feed-forward neural nets

- Definition

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Neural networks are non-linear estimators of the form

$$f(w, x) = \sigma \left(w_s^\top \sigma \left(\dots, \sigma(w_1^\top x) \right) \right)$$

where for all $i \in \{1, \dots, s\}$, $w_i \in \mathbb{R}^{n_i \times n_{i-1}}$

$\sigma : u \mapsto (\sigma(u_1), \dots, \sigma(u_p))^\top$ applies the scalar function σ to each coefficient of the input vector.

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- The integer s is often called the **depth** of the neural net. A neural network is called a deep neural net as soon as $s \geq 2$

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Neural networks are non-linear estimators of the form

$$f(w, x) = \sigma \left(w_s^\top \sigma \left(\dots, \sigma(w_1^\top x) \right) \right)$$

where for all $i \in \{1, \dots, s\}$, $w_i \in \mathbb{R}^{n_i \times n_{i-1}}$

$\sigma : u \mapsto (\sigma(u_1), \dots, \sigma(u_p))^\top$ applies the scalar function σ to each coefficient of the input vector.

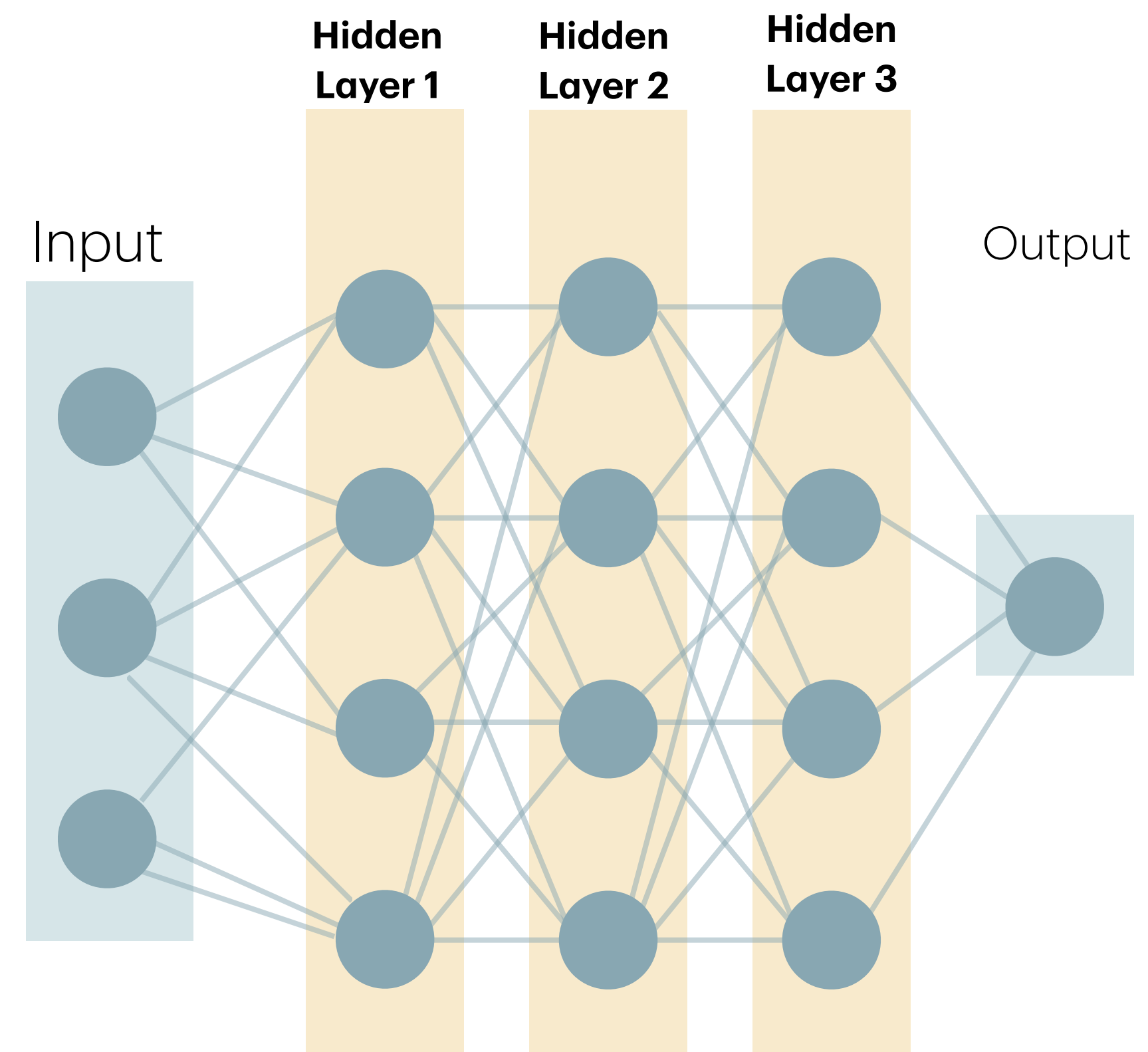
Remarks :

- n_0 must match the dimension d of the input x .
- σ is often called an **activation function**.
- Examples of standard activation functions include

$$\text{ReLU} : t \mapsto \max(t, 0) \quad \text{Sigmoid} : t \mapsto \frac{1}{1 + \exp(-t)} \quad \text{Tanh} : t \mapsto \frac{e^t - e^{-t}}{e^t + e^{-t}}$$

- The integer s is often called the **depth** of the neural net. A neural network is called a deep neural net as soon as $s \geq 2$

■ Visualisation



A complexity result for feed-forward neural nets

■ A recursive bound

Proposition : A recursive bound

Let $\mathcal{L}$ be a class of functions from $\mathbb{R}^d$ to $\mathbb{R}$ that includes the zero function.

Let $\sigma : \mathbb{R} \rightarrow \mathbb{R}$ be γ -Lipschitz.

Define

$$\mathcal{L}' := \left\{ x \in \mathbb{R}^d \rightarrow \sigma \left(\sum_{j=1}^m w_j l_j(x) + b \right) \in \mathbb{R} : |b| \leq \beta, \|w\|_1 \leq \omega, l_1, \dots, l_m \in \mathcal{L} \right\}.$$

Then, for $x = (x_1, \dots, x_n) \in \mathcal{X}^n$,

$$\text{Rad}(\mathcal{L}'_x) \leq \gamma \left(\frac{\beta}{\sqrt{n}} + 2\omega \text{Rad}(\mathcal{L}_x) \right),$$

where

$$\mathcal{L} = \{\ell(x_1), \dots, \ell(x_n), \ell \in \mathcal{L}\},$$

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Proof : on the black board.

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Proof : on the black board.

■ Complexity result

Proposition : A recursive bound

Let σ be γ -Lipschitz.

Consider the class $\mathcal{L}$ of neural networks losses with n layers

$$\ell^{(k)}(x) := \begin{cases} \sigma(\mathbf{w}^{(k)}x + b^{(k)}) & \text{if } k \leq n-1 \\ \mathbf{w}^{(n)}x + b^{(n)} & \text{if } k = n \end{cases} \quad \text{such that}$$

$$\|w^{(k)}\|_{\infty} \leq \omega, b^{(k)} \leq \beta, \forall k \in \{1, \dots, n\}.$$

Then, for any $x = (x_1, \dots, x_n) \in \mathbb{R}^{d \times n}$, the set $\mathcal{L}_x$ satisfies

$$\text{Rad}(\mathcal{L}_x) \leq \frac{1}{\sqrt{n}} \left(\beta + 2\omega\beta\lambda \sum_{k=0}^{\iota-3} (2\omega\lambda)^k + 2\omega(2\omega\lambda)^{\iota-2} \max_i \|x_i\|_{\infty} \sqrt{2 \log(2d)} \right).$$

Proof : exercise!